

UNIVERSITY COLLEGE LONDON

EXAMINATION FOR INTERNAL STUDENTS

MODULE CODE : MATH3801

**ASSESSMENT : MATH3801B
PATTERN**

MODULE NAME : Logic

DATE : 09-May-14

TIME : 10:00

TIME ALLOWED : 2 Hours 0 Minutes

All questions may be attempted, but only marks obtained on the best four solutions will count.

The use of an electronic calculator is not permitted in this examination.

1. (a) Define \mathcal{L} , the set of formulae in the general first order predicate language.
 - (b) For each of the following strings, use the definition of \mathcal{L} in order to determine whether or not the string is a formula:
(below, x, y, z are variable symbols and P is a unary predicate symbol)
 - (i) $\forall x \Rightarrow Px \Rightarrow \neg PyPz$
 - (ii) $\forall x \Rightarrow Px\forall yPy\neg Pz$
 - (c) Convert each of the following formulae of \mathcal{L}_{math} to a formula in \mathcal{L} :
(below, x, y, z are variable symbols and P is a unary predicate symbol)
 - (i) $((Px \vee Py) \vee Pz)$
 - (ii) $(\exists y)((Px \Rightarrow Py) \wedge Py)$
 - (d) Define the weight of a string.
 - (e) Show that any formula has weight -1 .
2. (a) Give the definition of a valuation on the set \mathcal{L}_0 of propositions.
 - (b) Suppose that S is a subset of the set \mathcal{L}_0 of propositions and that α is a proposition. State what it means to say that the semantic implication $S \models \alpha$ holds.
 - (c) Use the semantic tableaux method to determine whether or not each of the following semantic implications holds (where α, β, γ are distinct primitive propositions). If a semantic implication does not hold, describe a (type of) valuation for which it fails to hold.
 - (i) $\{\alpha \Rightarrow (\beta \Rightarrow \gamma)\} \models (\beta \Rightarrow (\alpha \Rightarrow \gamma))$
 - (ii) $\{\alpha \Rightarrow (\beta \Rightarrow \gamma)\} \models (\gamma \Rightarrow (\beta \Rightarrow \alpha))$
 - (iii) $\{(\neg\alpha) \wedge (\beta \Rightarrow \gamma), (\neg\beta) \Rightarrow \alpha\} \models (\beta \wedge \gamma)$

3. (a) Give a direct proof of the following (for propositions α, β):

$$\{\alpha \Rightarrow (\alpha \Rightarrow (\neg\beta))\} \vdash (\alpha \Rightarrow (\neg\beta))$$
 (You may assume the validity of the following theorem: $\vdash (\alpha \Rightarrow \alpha)$.)
- (b) (i) State the Deduction Theorem for propositional logic.
 (ii) Use the Deduction Theorem to show that (for propositions $\alpha, \beta, \gamma, \delta$):

$$\{\alpha \Rightarrow ((\beta \Rightarrow \gamma) \Rightarrow \delta)\} \vdash \alpha \Rightarrow (\gamma \Rightarrow \delta)$$
- (c) (i) State what it means for a set S of propositions to be consistent.
 (ii) State and prove the Adequacy Theorem for propositional logic.
 (You may assume that any consistent set of propositions has a model.)

4. (a) State and prove the Compactness Theorem for first order predicate logic.
 (You may assume the following form of the Completeness Theorem for first order predicate logic: If S is a set of sentences in a first order predicate language, then S is consistent if and only if S has a model.)
- (b) Describe a theory in a suitably defined first order predicate language $\mathcal{L}(\Pi, \Omega)$, such that a structure U is a (normal) model of the theory if and only if U is a poset (with equality) containing at most 5 elements.
- (c) Does there exist a theory in a suitably defined first order predicate language $\mathcal{L}(\Pi, \Omega)$, such that a structure U is a (normal) model of the theory if and only if U is a poset (with equality) containing finitely many elements?
 Explain your answer; you may quote, without proof, any relevant results.

5. (a) Define the notion of a register machine, giving also a description of what a program is and of the types of instructions associated to the states of a program.
- (b) Describe a way that may be used to uniquely encode each program, by a natural number.
- (c) Consider the following function:

$$f : \mathbb{N}_0 \rightarrow \mathbb{N}_0, \quad f(m) = 4$$

- (i) Describe a register machine that shows that f is computable.
 (ii) Find a natural number that uniquely encodes the described program.
- (d) Consider the following function:

$$g : \mathbb{N}_0 \rightarrow \mathbb{N}_0, \quad g(m) = 3m + 2$$

Describe a register machine that shows that g is computable.